CSC363 Tutorial 10 This will take like 40 mins, i'm guessing?

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University of Chungus in Japanese dub

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## Learning objectives this tutorial

By the end of this tutorial, you should...

Have an intuitive understanding of what co-NP is.

Be able to define what NP-hard is, and appreciate proofs of NP-hardness, because ironically proving NP-hardness is easier than proving NP-completeness.

Be a master of the famous "educational" game *Kahoot*, which is much more enjoyable than amogus.

Big Chungus certified readings: what readings? Iol. the lecture hasn't even happened yet.

#### ААААААААААААААААААААААААААААААА



Mohammad Mahmoud Tue 23/03/2021 13:29

To: Eric Lauw; Yousef Akiba <yousefakiba@gmail.com>; Muhammad Huzaifa; Daniel Ceniceros; Paul Zhang

Also for this week tutorial, introduce any example you see helpful, and perhaps introduce co-NP

oh no... not again ;-;

so uh, i hope youse like kahoot! i dunno what else i could cover... : $(^1$ 

 $<sup>^1</sup>$ last year when i was taking CSC363 and everything moved online, tutorials were just straight up cancelled! D:



mmm... I briefly mentioned this definition verbally in the last tutorial, but I don't know if any of you remember D:

Recall (maybe): A language A is NP-complete if

(a) 
$$A \in NP$$
;

(b) For every  $B \in NP$ ,  $B \leq_{p} A$ .

Now, sometimes we get lazy in proving some problem is NP-complete and forget to prove (a) are unable to prove (a) easily. You'll see some examples later. Despite that, we might still be able to prove (b)!

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### NP-hard 😳

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**Task:** speedrun proving NP-complete  $\subseteq$  NP-hard fullmarks%.





Intuitively, NP-hard problems are "at least as hard as the hardest problems in NP".

Here are some NP-hard problems:

Every NP-complete problem is NP-hard.

The **Travelling Salesman Problem**<sup>2</sup> is NP-hard; we don't know if it's in NP or not.

Super Mario Brothers is NP-Hard.

The Halting Problem

 $HP = \{(M, w) : M \text{ is a TM that halts on } w\}$ 

is NP-hard, but is known to not be NP, because it isn't even computable!

 $^2$ lf you haven't heard of it before, search it up! It's useful for food delivery n stuff maybe

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Since NP is the set of languages with poly-time verifiers, a language A is in co-NP if and only if there is a poly-time verifier V for its complement:

 $\overline{A} = \{w : \text{there exists } c \text{ such that } V(w, c) \text{ accepts}\},\$ 

or equivalently,

 $A = \{w : \text{there is no } c \text{ such that } V(w, c) \text{ accepts} \}.$ 

This means we can *verify that something is not in our language* in polynomial time.

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**Task:** Answer the following question: is NP = co-NP?

Actually, we don't know if NP = co-NP. But the following problems are co-NP!

pretty much any problem in NP, if you just take its complement. e.g. turn SAT (which asks "is this formula satisfiable?") into the problem "is this formula unsatisfiable?"

NOTE: "co-NP" is not the same as "not NP"! Be careful.

## yahoot time

(gotta fill time somehow, there isn't much to cover today though!)

winner gets (imaginary) \$10 sushi juice coupons. when i open a sushi juice store, you will be able to redeem those coupons.

